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08 – Support to Mobile Messaging and Event Management

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Decoupled Communications: Messaging

As already stated, **relevance of decoupling** in communication and interaction among mobile distributed components

Sometimes **message exchange** is even used as the general term to indicate the primary type of **mobile communication middleware** (see S. Tarkoma, “Mobile Middleware”) to highlight the importance of decoupling

Any mobile messaging solution must define:

- Principles and architecture
- Message **syntax**
- Protocol** for message exchange
- Locator**

Sometimes protocol term is used to include also syntax and locator...



Messaging: Principles and Architecture

Primary principle: **loose coupling** (via **standard and open protocols/formats**)

In real systems, also **extensibility**. How to?

- ❑ **Version identifiers included in messages** (non-recognized versions are considered as errors; back-compatibility?)
- ❑ **Formats with extension points**
- ❑ **Forward compatibility with possibility to ignore** message parts that are not recognized (example of application of **robustness principle**)

Usually **middleware APIs** to allow applications to use communication facilities; sometimes **middleware with visibility** of requirements for data exchange and their semantics



Messaging: Message Syntax

Marshalling/unmarshalling management:

- ❑ **Implemented at the application level**
- ❑ **Code may be automatically generated** (typically based on approaches like Interface Description Language – IDL – which is considered at development time)
- ❑ **Introspection** (higher expressive power for developers but typically more expensive at runtime)

How to agree on data format?

- ❑ **Specification** (usual approach of Internet protocols with messages in binary format)
- ❑ **Negotiation**
- ❑ **Receiver-makes-right** (sender uses its native formats and specifies metadata to indicate which formats are used)

Primary types of message formats:

- ❑ Binary (ASN.1, ...) or text-based (XML, JSON, ...)



In addition to the usual protocol properties for communication in distributed systems (headers with metadata and payload, also application-layer metadata, message types and with which exchange patterns, ...), special accent on **connection management**

- ❑ Protocols that “mimic” transport layer, with application-level connection in 1:1 relationship with transport-level connection
- ❑ **More often protocols that decouple the two aspects (persistent session feature over multiple and temporary transport connections**; see TCP and change of dynamic IP address, or SIP...)

«Pure» end-to-end perspective or usage of mediators up to the application level?

Wide usage of **store-and-forward architectures and protocols** (decoupling in time, optimization of implementations for reliability, violation of end-to-end principle)



Classical schemes for message exchange: one-way, two-way, request-response, subscribe-notify, conversation, ...

Relevant:

- ❑ Role at transport layer (**initiator-listener**), not necessarily the same as for the application/messaging levels
- ❑ Usual distinction **blocking vs. non-blocking**
 - **Polling method**, usually with object (**promise or future**) given to the sender; possible to make either inspect or blocking claim
 - **Callback**

Which of the two schemes facilitates more the development of mobile applications and/or for mobile systems?



Messaging: Locator

We are used to **locators strongly tightened to network addresses**

But also locators more articulated and complex, e.g., which include port numbers (transport) or paths (URLs)

In mobile systems **many types of locators**, also non IP-based, in particular in the past when IP was not so dominant

Anyway, even nowadays, possibility of:

- ❑ **“Transparent” locators**, typically implemented as URIs and codified in XML (it increases the level of abstraction + decoupling)
- ❑ **“Opaque” locators**, as in CORBA. Need of middleware to generate and use opaque locators

Is mobility management a network-layer issue? Of course, given what we have already widely discussed in the course, NOT ONLY...

Often it is written that *mobile hosts are managed as second-class citizens; towards locators independent from network layer...*



Messaging: Design for Mobility

Usual general considerations on:

- ❑ **Valuable role of proxies**, e.g., to split transport connections in two parts (breaking the end-to-end principle)
- ❑ **Problems of Network Address Translation** (NAT) when mobile nodes are willing to offer services (see also FTP and IRC...)

In addition:

- ❑ **Ability to complete the scheme of message exchange** even if communicating entities move
- ❑ Exploitation of classic transport-level connections is usually preferred
- ❑ **Simple syntax and reduced message content**, also considering compression facilities
- ❑ **Security at message level**: with SSL-like approaches (connection-level security), which kind of issues if end-to-end principle is broken? **Message-level security with security applied to (parts of) the payload, not to headers**; of course, also combination of connection and message



As usual:

- ❑ Distinction between ***end-to-end and hop-by-hop***
- ❑ Basic technique with ***acknowledgement and re-transmissions*** (also at the application level)

ACK types:

- ❑ Regular
- ❑ Cumulative
- ❑ Negative
- ❑ Piggy-backed

In-order delivery? Sometimes it can be sacrificed for efficiency motivations

Indeed, reliability reduction due to performance motivations is a well-known concept (DNS, NTP, SIP, ... typically use UDP)



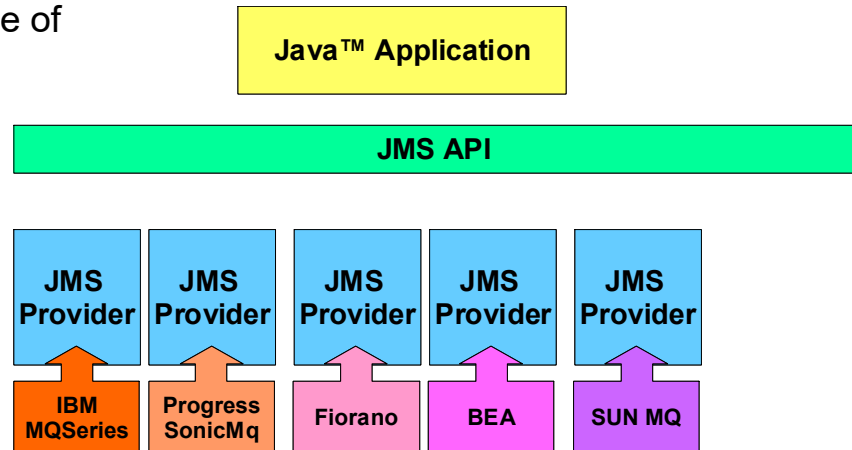
- ❑ Possibility to ask for ***only-once semantics for message delivery*** (more precisely once-and-only-once for persistent usage, at-most-once for non-persistent usage)
- ❑ ***Decoupling in time thanks to durable destinations***
- ❑ Partial time coupling for topics: it can be reduced via ***durable subscriptions***
- ❑ Possibility of ***non-blocking reception via listeners***
- ❑ Usage of ***ConnectionFactory***
- ❑ Messages sent within a session (via a given Session object) towards the same destination ***benefit from in-order delivery property***
- ❑ ***Three types of ACK*** (lazy, automatic, and client-side)



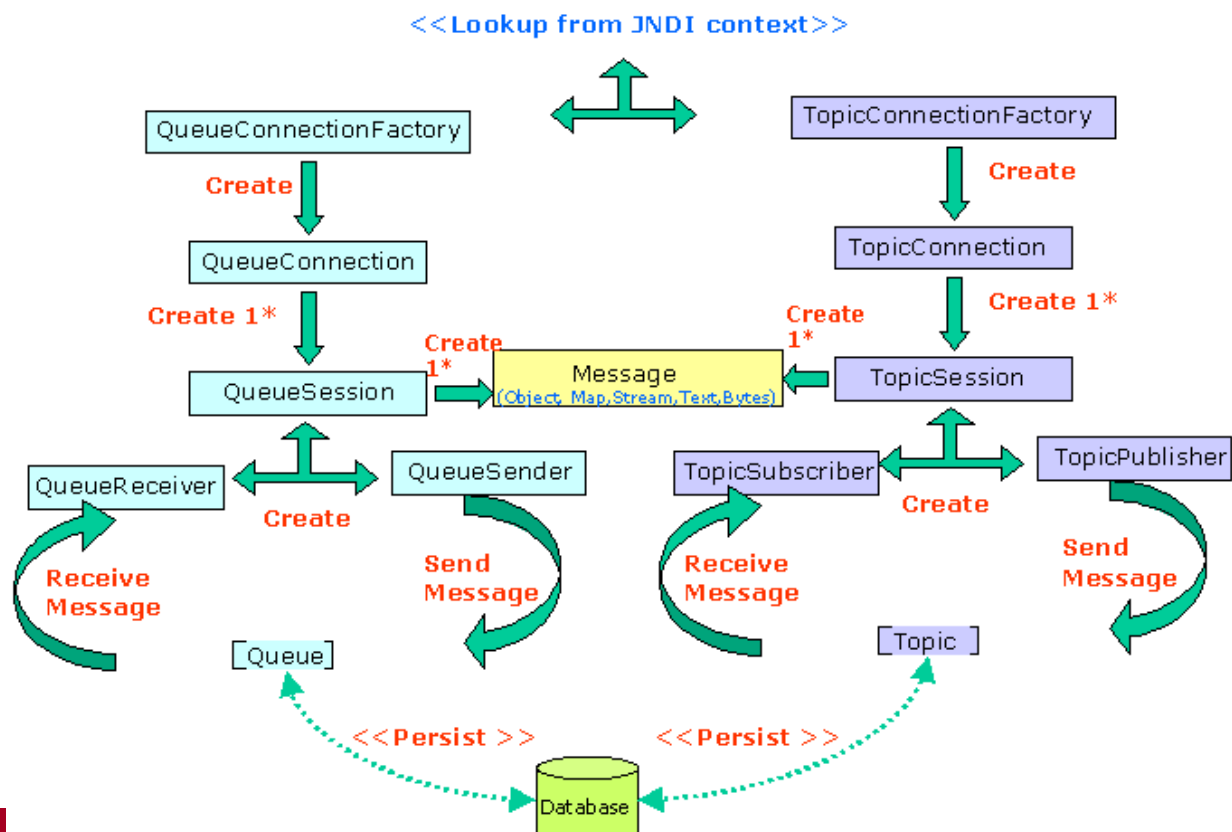
Design Goals in JMS

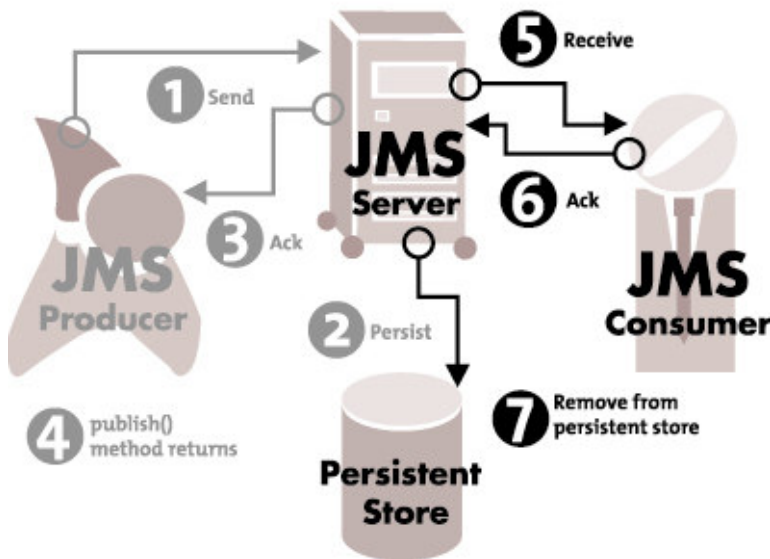
JMS is part of the J2EE platform. Goals:

- ❑ **Compliance/similarity** with APIs of **existing messaging systems**
- ❑ **Independency** from vendors of messaging systems
- ❑ Coverage of most common facilities that are offered in messaging systems
- ❑ It promotes the usage of Java technology



“Graphical Summary” of JMS APIs





- ❑ **Producer-side and consumer-side perspectives**
- ❑ Differences between **persistent and non-persistent cases**
- ❑ When is it possible to have **duplicated messages**?
- ❑ When is it possible to have **message losses**?
- ❑ In addition, **three differentiated types of ack**

- ❑ **PERSISTENT**
 - Default
 - It specifies to **JMS provider to guarantee** that the message **is not lost when in transit**, e.g., because of a failure at the JMS provider
- ❑ **NON_PERSISTENT**
 - **It does NOT request storing messages** at the JMS provider side
 - Better performance results

[SetDeliveryMode\(\)](#) method in the [MessageProducer](#) interface

- `producer.setDeliveryMode(DeliveryMode.NON_PERSISTENT);`
- **Extended form:** `producer.send(message, DeliveryMode.NON_PERSISTENT, 3, 10000);`



Priority and Expiration in Message Delivery

- 10 priority levels
 - from 0 (lowest) to 9 (highest)
 - default = 4

Usage of `setPriority()` method of `MessageProducer` interface, e.g., `producer.setPriority(7)` ;

or the extended form `producer.send(message, DeliveryMode.NON_PERSISTENT, 7, 10000)` ;

- Expiration: possibility to **configure TTL** via `setTimeToLive()` of the `MessageProducer` interface
 - `producer.setTimeToLive(60000)` ;
 - Or extended form, `producer.send(message, DeliveryMode.NON_PERSISTENT, 3, 60000)` ;



Messaging Examples: CORBA Messaging

CORBA Messaging specification includes:

- Asynchronous Messaging Interface (AMI)
 - Possibility of both **polling and callback** (callback is passed as CORBA object, therefore even not in the same addressing space of client)
- Time Independent Invocation (TII) to specify which CORBA objects play the **role of router** for the message
 - Rationale: sender and recipient may be temporarily disconnected
 - They compose a **store-and-forward network**

CORBA locator = Interoperable Object Reference (IOR), with different profiles depending on binding protocol

Messages in binary format = Common Data Representation (CDR)

Extreme flexibility in the choice of the protocol



CORBA AMI: Callback Mode

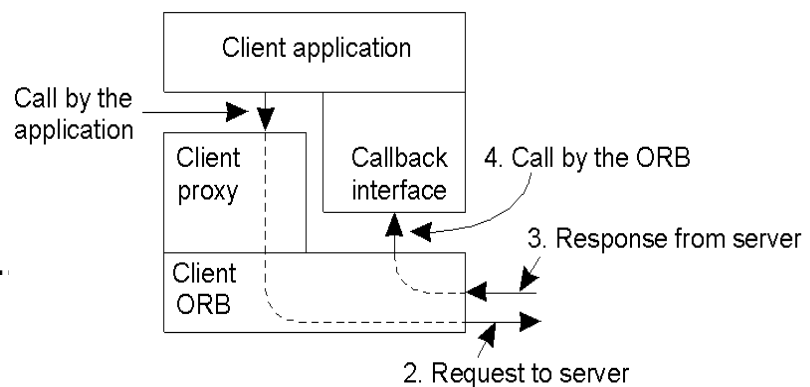
Callback: client provides callback method to be invoked by the support after service completion via a given **fire-and-forget (automatically invoked)**

In place of: `int somma (in int i, in int j, out int somma)`
`void sendcallback_somma (in int i, in int j)`
`void callback_somma (in int success, in int somma)`

Usage of two methods
by changing **only**
client implementation
and **NOT any**
service part

Client invokes **sendcallb...**

ORB invokes **callback_som...**

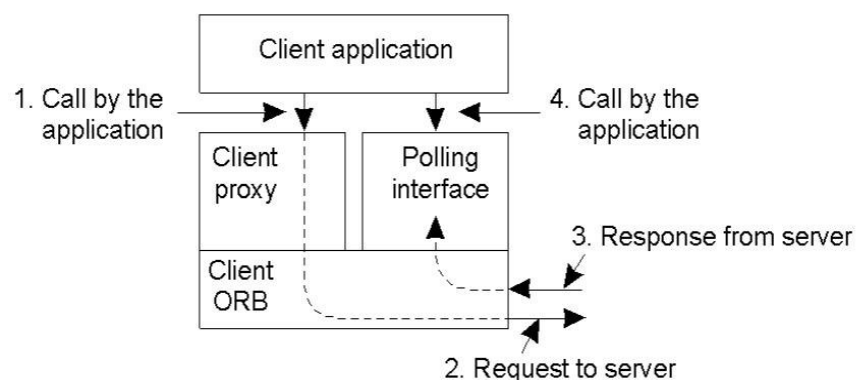


CORBA AMI: Polling Mode

Asynchronous polling: client decides **when and whether** to interrogate a method to check completion of remote operation (by collecting results); this method is **created by the messaging support**

In place of: `int somma (in int i, in int j, out int somma)`
`void sendpoll_somma (in int i, in int j)`
`void pollsomma (out int success, out int somma)`

Result is collected on request by invoking **pollsomma operation that is autom. generated by CORBA support**





Messaging Examples: Extensible Messaging and Presence Protocol (XMPP)

Essentially designed for instant messaging

RFC 3920 is oriented and similar to the existing implementation of the Jabber protocol; good popularity and widespread utilization thanks to the adoption by *Google, Twitter, Facebook, ...*

It includes ***publish/subscribe mechanisms*** (see the following slides...) ***to update presence and state, and for service discovery***

Client-server model: client sends an XMPP dataflow to a server, after parameter negotiation

Peer-to-peer model: servers coordinate together for delivery to recipients

Usage of so-called ***stanzas***, of three types:

- ❑ *Message stanza* – one-to-one communication, similar to emails
- ❑ *Presence stanza* – simple pub/sub mechanism, communication is transferred to all subscribers
- ❑ *Info/Query stanza* – request-response mechanism



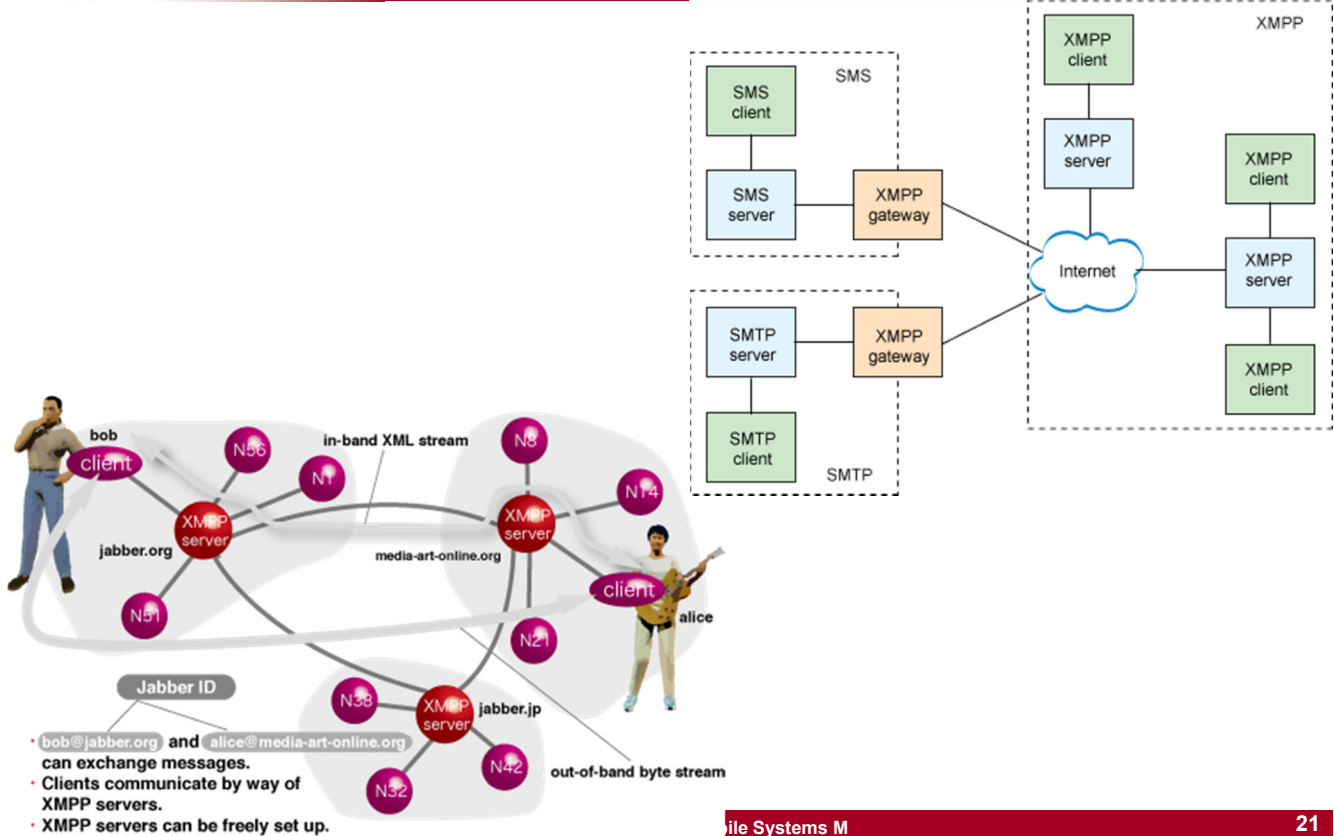
Messaging Examples: Extensible Messaging and Presence Protocol (XMPP)

XMPP messages are streams codified in XML

Given the widespread adoption, good candidate to support messaging in mobile systems, EVEN IF:

- ❑ Not specifically designed for mobile systems
- ❑ Expensive XML processing, expensive connection management in particular in terms of energy
- ❑ ***Expensive re-connections to XMPP server*** (need to re-establish a ***new session of interaction per any new transport connection***, transmission of XML data that is non-negligible at each session start)

Android implements a specialized and proprietary variant of it, with non-XML-based protocol and NO creation of a new session per any new connection



SOAP is built on top of interaction model based on message exchange

- ❑ Architecture based on senders, receivers, and intermediate nodes
- ❑ Locator = HTTP URI
- ❑ **Document-style SOAP: messages as XML-based documents** that have to be processed
- ❑ Possibility of **different protocol bindings**, but definitely the most used one is HTTP, utilization of POST method (employed more as transport protocol, while ignoring its application semantics)
- ❑ In mobile environments, where HTTP is sometimes the only protocol practically usable because of firewalls and NAT, this use/misuse of HTTP could be considered as legitimate and become largely adopted...
- ❑ **Also specification for binding to email and XMPP**



Messaging Examples: Representational State Transfer (REST)

REST is substantially a ***solution architectural style***,
Resource Oriented Architecture (Roy Fielding, UCI PhD
Thesis, 2000)

To promote ***client-server and stateless interaction***,
oriented to the usage of caching opportunities, also
with possibility of code-on-demand to clients

***Any resource has a persistent identifier; idea to
transfer NOT resources but their representations via
HTTP protocol***

Constraint: exchange of self-descriptive messages (languages for
representation, negotiation of supported modes, ...)



Messaging Examples: Representational State Transfer (REST)

Locator = HTTP URI

Three types of metadata included in HTTP headers:

- ❑ ***Resource metadata*** – about resources, e.g., timestamp about last modification
- ❑ ***Representation metadata*** – about transferred representation, e.g., its media type
- ❑ ***Control metadata*** – about message, e.g., its length and caching possibility

Notable example: ***RESTful Web services***

RESTful Web service as a simple Web service implemented by using HTTP
and REST principles, thus resource collection with 3 well-defined
aspects:

- ❑ ***URI base for service***, e.g., <http://example.com/resources/>
- ❑ ***Internet media type*** for data used in the service (usually JSON or XML)
- ❑ Set of service operations supported ***via HTTP*** method invocations (e.g.,
via POST, GET, PUT or DELETE)



Messaging Examples: Representational State Transfer (REST)

Notable example: *RESTful Web services*

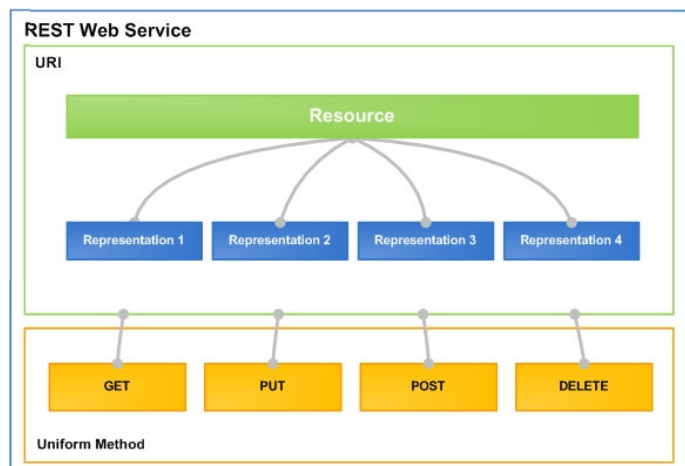
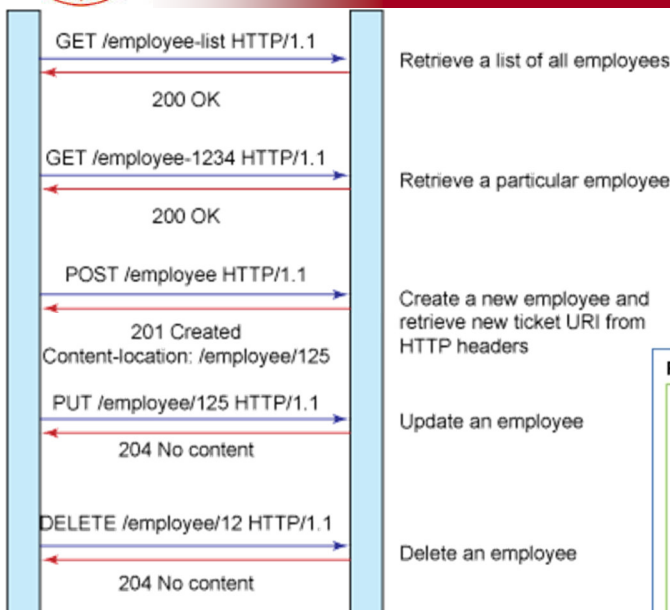
Risorsa	GET	PUT	POST	DELETE
URI for resource collection, e.g., http://example.com/resources/	To list all collection members	To replace the whole collection	To create a new element to be inserted in the collection	To remove the whole collection
URI for single element, e.g., http://example.com/resources/ef7d-xj36p	To obtain the representation of the targeted element, expressed in the appropriate Internet media type	To replace or create an element of the collection	To consider the element as a collection and to create a new element internally to it	To remove an element from the collection

Examples of today's REST usage:

- ❑ Majority of **Web blogs** (download of XML files in *RSS/Atom format*, which contain links to other resources)
- ❑ **Simple Storage Service (S3)** by Amazon.com
- ❑ **OpenStreetMap** (REST interface)... and many many others



Messaging Examples: Representational State Transfer (REST)



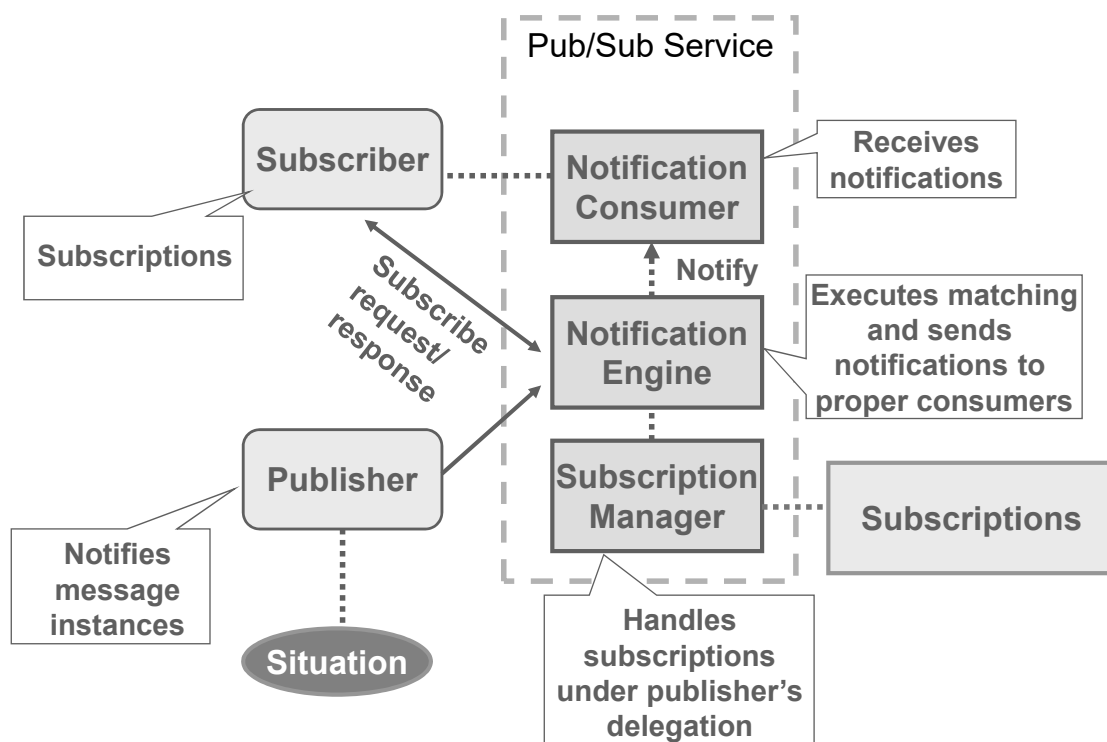


Event Management and Publish/Subscribe Systems

- ❑ **Event delivery from publishers to subscribers**
 - Events as messages with content
 - **One-to-many, many-to-many** (traditional message systems are queue-based and one-to-one)
 - Often implemented based on **messaging systems and on store-and-forward solutions**
- ❑ Comm. paradigm of frequent usage, in particular in mobile systems
 - **Decoupling in space and time**
- ❑ Event system as **logically centralized system**
 - Anonymous communication
 - Possibility to **use filters** (on headers or entire messages)
 - Basic primitives: subscribe, unsubscribe, publish, also with filters
- ❑ **Different topologies for routing and different semantics** associated to event sending/notification
- ❑ Associated operations are typically **non-blocking** (polling, callback)



General Architecture for Publish/Subscribe Systems





Event Router and Topologies

Event router or broker

- ❑ Works as mediator (decoupling) between publishers and subscribers
- ❑ Usage of **routing table (also with filters)** for local event dispatching or to indicate to which «near» router to forward in the case of distributed brokers (to achieve scalability, reliability, and high availability)
- ❑ **Filters** may be also based on **content** => **content-based routing**
- ❑ Other non-functional requirements: notification within time deadlines (**bounded delivery time**), QoS, fault-tolerance, ordering (**causal order, total order**)

Possible router topologies:

- ❑ **Centralized**
- ❑ **Hierarchical** (notifications always sent to master, i.e., root of the distribution tree)
- ❑ **Cyclic, acyclic** (peer-to-peer, cyclic allows redundancy but need of *minimum spanning tree techniques* to prevent from cycles)
- ❑ Based on **rendez-vous point** (special router that works as rendez-vous, typically for pre-determined types of events)

Partially related: have you ever heard of **Distributed Hash Tables (DHTs)**?



Interest Propagation and Subscriptions

One of the primary functions of a router is to **propagate notifications to near routers that are interested** in that event. To this purpose, how to propagate interests and subscriptions?

Properties to be achieved: reduced forwarding overhead, high performance, fast support to variations

- ❑ **Simple routing**: any router knows all subscriptions in the global systems (subscription flooding), possibly with optimization of NO forwarding if subscription message has been already circulated
- ❑ **Covering-based routing**: forwarding of only the more general subscription filters (*which possible issues with unsubscription?*)
- ❑ **Merging-based routing**: it allows to merge different entries in routing table for the sake of table size optimization (usually combined with covering, here also unsubscription issues)

Notifications are usually distributed over **reverse paths (wrt subscription paths)**



Decision about Message Routing

Depending on what is used to take message routing decisions, classification into:

- ❑ **Channel/topic-based**: depending on the channel (usually named channel) on which the event is published. Pub/sub agreement on the channel name, also possibility of associated multicast address
- ❑ **Subject-based**: depending on event subject, single field of info
- ❑ **Header-based**: depending on a set of fields. For example, SOAP supports header-based routing for its messages
- ❑ **Content-based**: possibly depending on the whole message content. Higher expressive power, higher costs

Also **context+content-based routing**, particularly suitable for mobile systems/services with **event filters that are context-dependent**



Java Model for Distributed Events

Also Java has a built-in **model for event distribution, based on RMI**, e.g., used in Jini/River

- ❑ Based on **Remote Event Listener** (consumers are registered to receive given types of events from given objects, `notify()` method)
- ❑ **Remote Event object returned back during notification** (data, reference to source object, handback object, unique identifier)
- ❑ **Lease mechanism**
- ❑ The specification includes possibility to define **Distributed Event Adaptors that implement filters and QoS policies**
 - Idea to exploit handback object, returned by the event source, to transfer state and behavior (e.g., to implement event filters)



Java Model for Distributed Events

```
package net.jini.core.event;

public class RemoteEvent {
    public long getID();
    public long getSequenceNumber();
    public java.rmi.MarshalledObject getRegistrationObject();
}
```

Events generated in local components may transfer even quite complex object state. **NOT distributed events: only info on how state retrieval is possible at runtime**

- ❑ Remote event as serializable object that can be transferred between listeners
- ❑ Idea, “stolen” from Xt Intrinsic and Motif solutions: **to register clients by including handback objects, returned back with any event**

For example, a Jini taxi driver subscribes to taxi bookings while passing through a city area (handback includes location); when it receives an event, it can be informed of old location (at the moment of registration)

Possibility to register other objects for notification delegation: in this case, handback can work as “reminder” with info of subscribers (*stock broker model*)



Java Model for Distributed Events

Event registration

Jini/River does NOT specify how to register listeners at event sources; only specification to use a class as return value from subscription:

```
package net.jini.core.event;
import net.jini.core.lease.Lease;

public class EventRegistration implements java.io.Serializable {
    public EventRegistration(long eventID, Object source,
                            Lease lease, long seqNum);

    public long getID();
    public Object getSource();
    public Lease getLease();
    public long getSequenceNumber();
}
```

Therefore, the developer of event source has to implement:

```
public EventRegistration
    addRemoteEventListener(RemoteEventListener listener);
```



Java Model for Distributed Events

Java model for local events work with objects that are all in the same addressing space

Jini as community of distributed objects that cooperate through **proxies**

For remote events, “***inversion of proxy direction***”

- ❑ For example, ***Jini client uses its proxy for service access and through it registers itself as listener***
 - Need for a proxy method to ***add event listeners***
- ❑ Proxy will invoke the “real” method for listener adding over the discovered resource
- ❑ Invocation of ***registration of local event to proxy***; invocation of proxy for remote resource registration

As if real resource ***obtains a proxy for the client to use in the notification chain***



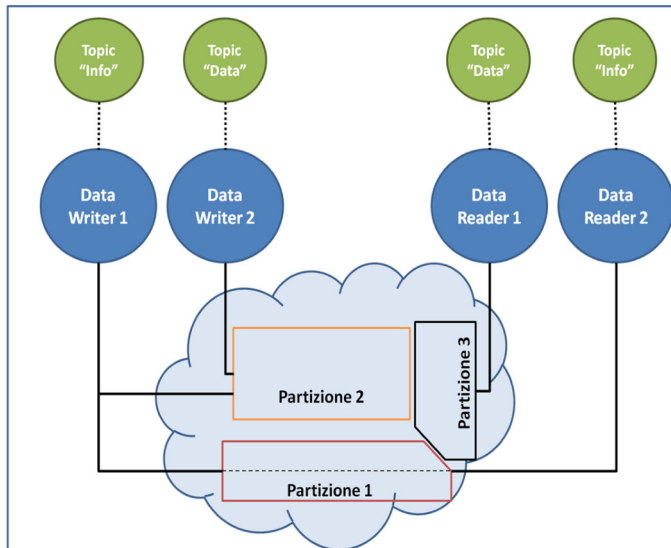
OMG Distributed Data Service (DDS)

- ❑ ***OMG specification*** (neither based on CORBA nor highly interoperable) for ***data distribution service designed for real-time systems***
- ❑ Specification defines ***APIs for so-called publish/subscribe data-centric communication***; in other terms, DDS middleware offers abstraction of global data space that is accessible to all interested applications
- ❑ Usage of ***combination of Topic objects and keys*** to univocally identify ***instances within a datastream*** of the same topic
- ❑ ***Support to content filtering and QoS negotiation***
- ❑ Suitable for distributed propagation of signals, data, and events

CORBA Event Service (NOT data-centric and with NO QoS support);
CORBA Notification Service (filters, QoS, but mandatory usage of CDR and IIOP)

Partitions are namespaces to allow the *logical splitting of a DDS domain*

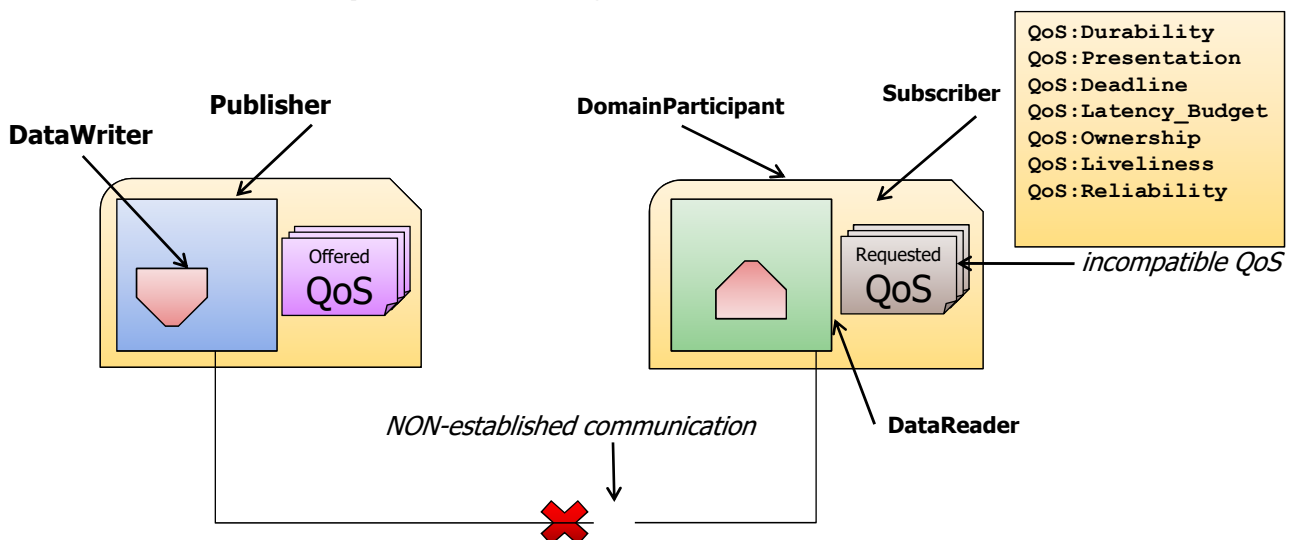
Publisher/Subscribers can decide *at runtime* (and NOT at instantiation time as for JMS Topics) on which partitions to publish/subscribe data



For a DataReader to receive messages from a DataWriter, there is the need to share **both the same Topic and the same partition**

Partitions are considered to enforce a QoS policy

- ❑ To allow a Subscriber receiving publications from a Publisher, **QoS properties have to be compatible**
- ❑ Protocol of **Request/Offer negotiation**



DDS supports different modes for message sending (e.g., best-effort, reliable) and personalized management of data persistence



Quality as Reliability

DDS identifies **two QoS policies for message reliability**:

- ❑ **BEST_EFFORT** – **NOT guaranteed** that all messages are received, NOT guaranteed delivery order
- ❑ **RELIABLE** – guaranteed that all messages are received and delivery order. Via Publishers that **re-send** data to Subscribers if needed and via Subscribers that send reception **feedback (ack)**

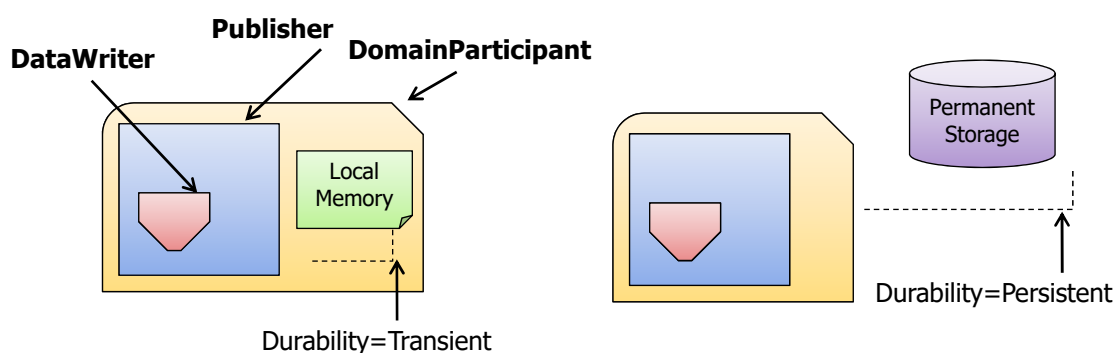
In reliable case, all sent messages are **kept in a history queue** while waiting for being confirmed (publisher side) and processed by application (subscriber side); queue size can be defined, through **HISTORY policy**

It is also possible to define **how many resources** (e.g., memory, max instances) to use to maintain data, through **RESOURCE_LIMITS policy**



Quality as Durability

- ❑ Through **Durability policy** it is possible to define whether and how many data to be maintained at publisher side in order to enable their future successive request
- ❑ DDS supports 3 persistency types:
 - **VOLATILE** – No Instance History Saved
 - **TRANSIENT** – History Saved in Local Memory
 - **PERSISTENT** – History Saved in Permanent storage





DDS Quality: Additional Policies

DDS supports a wide set of other policies to define:

- ❑ **Ordering of received messages** (DESTINATION_ORDER - BY_RECEPTION_TIMESTAMP, BY_SOURCE_TIMESTAMP – **eventual consistency, ...**)
- ❑ **Message priority** (LATENCY_BUDGET)
- ❑ Exclusiveness on some given data types (OWNERSHIP)
- ❑ Data authentication and security (USER_DATA)
- ❑ **Time constraints on message sending/delivery rates** (TIME_BASED_FILTER)
- ❑ Fault detection and heartbeat (LIVELINESS)

More detailed technical documents at :

- **Getting Started Guide**
www.rti.com/eval/rtidds44d/RTI_DDS_GettingStarted.pdf
- **RTI DDS User's Manual**
www.dre.vanderbilt.edu/~mxiong/tmp/backup/RTI_DDS_UsersManual.pdf



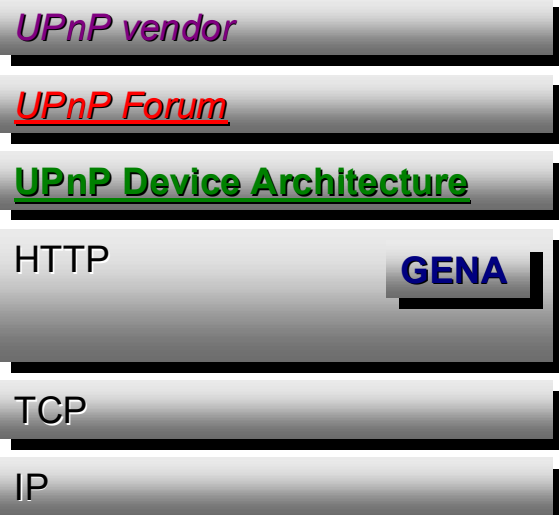
General Event Notification Architecture (GENA)

As already stated, used primarily in UPnP

- ❑ Control point is listener of modifications of device state
 - 0 obtains address
 - 1 discovers device
 - 2 determines XML descriptor
 - Obtains URL for eventing
 - 4 registers itself

Extreme simplicity:

Notification sending/reception via HTTP over TCP/IP or multicast UDP





GENA: Subscription

Control point has to register itself before being able to receive any event

```
SUBSCRIBE publisher path HTTP/1.1  
HOST: publisher host:publisher port  
CALLBACK: <delivery URL>  
NT: upnp:event  
TIMEOUT: Second-requested subscription duration
```

Device accepts subscription: it immediately sends a special event (initial) to control point with the value of all state variables

```
HTTP/1.1 200 OK  
SID: uuid:subscription-UUID  
TIMEOUT: Second-actual subscription duration
```



GENA: Notifications

When a state variable changes value at a device:

```
NOTIFY delivery path HTTP/1.1  
HOST: delivery host:delivery port  
CONTENT-TYPE: text/xml  
NT: upnp:event  
NTS: upnp:propchange  
SID: uuid:subscription-UUID  
SEQ: event key  
  
<e:propertyset xmlns:e="urn:schemas-upnp-org:event-1-0">  
  <e:property>  
    <variableName>new value</variableName>  
  </e:property>  
  Other (possible) names of variable and associated values  
</e:propertyset>
```



SIP Event Framework

Asynchronous notifications are essential to implement many SIP services (automated callback services, list of buddies who are currently online, message waiting services, ...)

To this purpose, **SIP Event Framework** (RFC 3265), based on SUBSCRIBE and NOTIFY methods

- ❑ SIP events are identified through three elements: Request URI, event type, and message body (optional)
- ❑ Notable example: **presence service, with so-called presentities and watchers**
 - Presence URI in the format `pres:paolo@domain`
 - **Scalability issues**
 - Need for **hierarchical organization** into domains and **action of even aggregation on localities** to reduce number of notifications



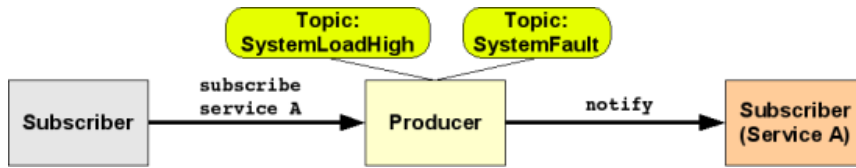
Web Services Event&Notification

Two key mechanisms to implement pub/sub for Web services: **WS-Eventing and WS-Notification** (standardization in 2006)

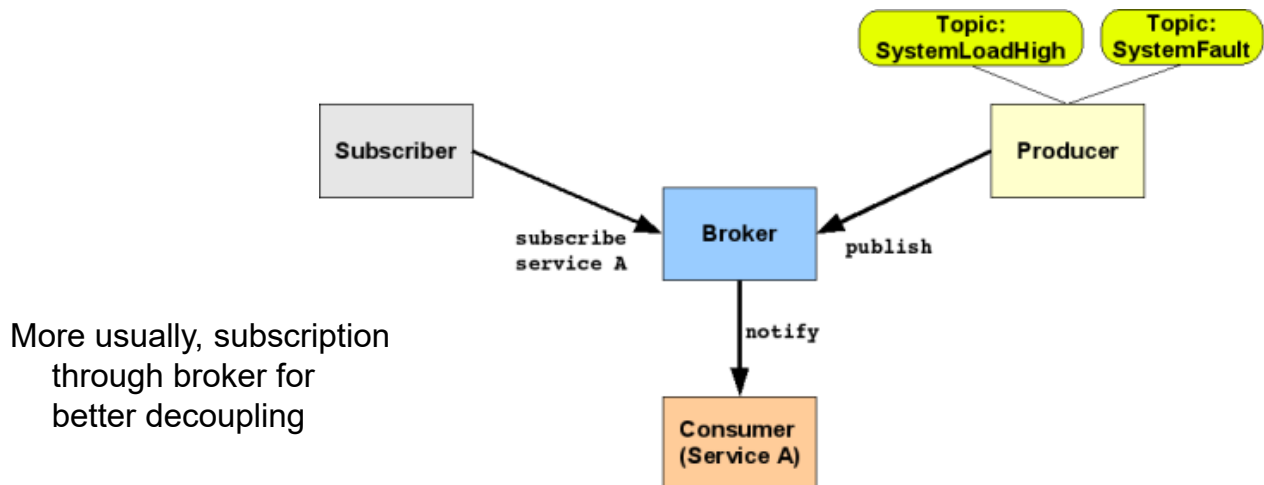
- ❑ **WS-Eventing** is the specification of protocol with which Web services have to **make/accept registrations for event notification**
 - Mechanisms to create/remove subscriptions
 - Mechanisms to define **expiration time** and to allow renewal
 - **Support to filters** (different languages for filter definition may be used)
- ❑ **WS-Notification** is the specification to allow Web services **to disseminate data to other Web services**
 - Also possibility of organizations oriented to interests (called topics) and **interest-based filtering**
 - **Distributed topologies for notification brokers**



Web Services Event&Notification



Possible subscription from third parties (direct, with NO broker)



More usually, subscription through broker for better decoupling



Programming Example of WS-Event&Notification

For example, *how to implement WS subscriber by using IBM WebSphere:*

- ❑ As usual, need to **obtain WSDL file** for notification broker and subscription manager services (resp. *NotificationBroker.wsdl* and *SubscriptionManager.wsdl*)
- ❑ If not yet available at client, need to execute wsimport to **generate client stub**
- ❑ **Look up at notification broker** (need for reference to notification broker service)
- ❑ **Instantiation of subscription request object and configuration of consumer reference**
- ❑ Instantiation of subscribe object to include subscription details, like reference to notification consumer

```
import org.oasis_open.docs.wsn.b_2.Subscribe;
import javax.xml.ws.wsaddressing.W3CEndpointReference;
import javax.xml.ws.wsaddressing.W3CEndpointReferenceBuilder;
// Crea oggetto subscription request. DEVE contenere
// ConsumerReference e PUO' includere filtro, InitialTerminationTime
// e SubscriptionPolicy
Subscribe subscribeRequest = new Subscribe();
W3CEndpointReference consumerReference = new
    W3CEndpointReferenceBuilder().address(consumerURI).build();
subscribeRequest.setConsumerReference(consumerReference);
```




Programming Example of WS-Event&Notification

Definition of *topic expression as registration filter*

It is possible to associate a *Filter object* to registration request to indicate which events are relevant (*filter based on topic, message content, or both*). For example, topic-based filter (with IBM helper classes):

```
import com.ibm.websphere.sib.wsn.jaxb.base.FilterType;
import com.ibm.websphere.sib.wsn.jaxb.base.TopicExpressionType;
// To prepare the topic expression
topicExpression = topicNamespacePrefix + ":" + topicExpression;
TopicExpressionType topicExpressionType = new TopicExpressionType();
topicExpressionType.setExpression(topicExpression);
// To specify mapping from namespace prefix to topic namespace URI
topicExpressionType.addPrefixMapping(topicNamespacePrefix,
    topicNamespace);
// To specify dialect TopicExpression to use
topicExpressionType.setDialect(topicDialect);
// Filter instantiation
FilterType filter = new FilterType();
// To add expression to filter and needed configuration
// subscribe with filter
filter.addTopicExpression(topicExpressionType);
subscribeRequest.setFilter(filter);
```



Programming Example of WS-Event&Notification

Specification of registration duration and request sending

Two modes to specify expiration time for registration:

- 1) namespace URI and QName objects
- 2) Helper factory = JAXB ObjectFactory

```
import javax.xml.bind.JAXBElement;
import javax.xml.datatype.DatatypeFactory;
import javax.xml.datatype.Duration;
// Option 1: Duration specification (one year from now)
DatatypeFactory factory = DatatypeFactory.newInstance();
Duration duration = factory.newDuration("1Y"; JAXBElement<String>
    initialTerminationTime = new JAXBElement<String>(
        new QName("http://docs.oasis-open.org/wsn/b-2",
            "InitialTerminationTime"), String.class, duration.toString());
// Option 2:
org.oasis_open.docs.wsn.b_2.ObjectFactory objectFactory = new org.
    oasis_open.docs.wsn.b_2.ObjectFactory();
initialTerminationTime = objectFactory.createSubscribeInitial-
    TerminationTime(duration.toString());

subscribeRequest.setInitialTerminationTime(initialTerminationTime);
org.oasis_open.docs.wsn.b_2.SubscribeResponse
    subscribeResponse = port.subscribe(subscribeRequest);
```