

Mobile Systems M

Alma Mater Studiorum – University of Bologna CdS Laurea Magistrale (MSc) in Computer Science Engineering

Mobile Systems M course (8 ECTS)

II Term – Academic Year 2019/2020

08 – Support to Mobile Messaging and Event Management

Paolo Bellavista paolo.bellavista qunibo.it

Luca Foschini luca.foschini@unibo.it

http://lia.disi.unibo.it/Courses/sm1920-info/



Decoupled Communications: Messaging

As already stated, *relevance of decoupling* in communication and interaction among mobile distributed components

Sometimes *message exchange* is even used as the general term to indicate the primary type of *mobile* communication middleware (see *S. Tarkoma, "Mobile Middleware"*) to highlight the importance of decoupling

Any mobile messaging solution must define:

- Principles and architecture
- Message syntax
- Protocol for message exchange
- Locator

Sometimes protocol term is used to include also syntax and locator...



Messaging: Principles and Architecture

Primary principle: *loose coupling* (via *standard and open protocols/formats*)

In real systems, also *extensibility*. How to?

- Version identifiers included in messages (non-recognized versions are considered as errors; back-compatibility?)
- Formats with extension points
- Forward compatibility with possibility to ignore message parts that are not recognized (example of application of robustness principle)

Usually *middleware APIs* to allow applications to use communication facilities; sometimes *middleware with visibility* of requirements for data exchange and their semantics



Messaging: Message Syntax

Marshalling/unmarshalling management:

- Implemented at the application level
- □ Code may be automatically generated (typically based on approaches like Interface Description Language IDL which is considered at development time)
- Introspection (higher expressive power for developers but typically more expensive at runtime)

How to agree on data format?

- Specification (usual approach of Internet protocols with messages in binary format)
- Negotiation
- Receiver-makes-right (sender uses its native formats and specifies metadata to indicate which formats are used)

Primary types of message formats:

□ Binary (ASN.1, ...) or text-based (XML, JSON, ...)



Messaging: Protocols

- In addition to the usual protocol properties for communication in distributed systems (headers with metadata and payload, also application-layer metadata, message types and with which exchange patterns, ...), special accent on *connection management*
- Protocols that "mimic" transport layer, with application-level connection in 1:1 relationship with transport-level connection
- More often protocols that decouple the two aspects (persistent session feature over multiple and temporary transport connections; see TCP and change of dynamic IP address, or SIP…)

«Pure» end-to-end perspective or usage of mediators up to the application level?

Wide usage of **store-and-forward architectures and protocols** (decoupling in time, optimization of implementations for reliability, violation of end-to-end principle)



Messaging: Protocols

Classical schemes for message exchange: one-way, two-way, request-response, subscribe-notify, conversation, ...

Relevant:

- Role at transport layer (*initiator-listener*), not necessarily the same as for the application/messaging levels
- Usual distinction blocking vs. non-blocking
 - > **Polling method**, usually with object (**promise or future**) given to the sender; possible to make either inspect or blocking claim
 - Callback

Which of the two schemes facilitates more the development of mobile applications and/or for mobile systems?



Messaging: Locator

We are used to *locators strongly tightened to network addresses*

But also locators more articulated and complex, e.g., which include port numbers (transport) or paths (URLs)

In mobile systems *many types of locators*, also non IP-based, in particular in the past when IP was not so dominant Anyway, even nowadays, possibility of:

- "Transparent" locators, typically implemented as URIs and codified in XML (it increases the level of abstraction + decoupling)
- "Opaque" locators, as in CORBA. Need of middleware to generate and use opaque locators

Is mobility management a network-layer issue? Of course, given what we have already widely discussed in the course, NOT ONLY...

Often it is written that mobile hosts are managed as second-class citizens; towards locators independent from network layer...



Messaging: Design for Mobility

Usual general considerations on:

- Valuable role of proxies, e.g., to split transport connections in two parts (breaking the end-to-end principle)
- Problems of Network Address Translation (NAT) when mobile nodes are willing to offer services (see also FTP and IRC...)

In addition:

- Ability to complete the scheme of message exchange even if communicating entities move
- Exploitation of classic transport-level connections is usually preferred
- Simple syntax and reduced message content, also considering compression facilities
- Security at message level: with SSL-like approaches (connection-level security), which kind of issues if end-to-end principle is broken? Message-level security with security applied to (parts of) the payload, not to headers; of course, also combination of connection and message



Messaging: Reliability

As usual:

- Distinction between end-to-end and hop-by-hop
- Basic technique with acknowledgement and retransmissions (also at the application level)

ACK types:

- Regular
- Cumulative
- Negative
- Piggy-backed

In-order delivery? Sometimes it can be sacrificed for efficiency motivations

Indeed, reliability reduction due to performance motivations is a well-known concept (DNS, NTP, SIP, ... typically use UDP)



Messaging Examples: Java Message Service (JMS)

- □ Possibility to ask for only-once semantics for message delivery (more precisely once-and-only-once for persistent usage, at-most-once for non-persistent usage)
- Decoupling in time thanks to durable destinations
- Partial time coupling for topics: it can be reduced via durable subscriptions
- Possibility of non-blocking reception via listeners
- Usage of ConnectionFactory
- Messages sent within a session (via a given Session object) towards the same destination benefit from in-order delivery property
- □ *Three types of ACK* (lazy, automatic, and client-side)



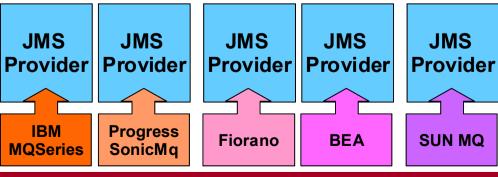
Design Goals in JMS

JMS is part of the J2EE platform. Goals:

- Compliance/similarity with APIs of existing messaging systems
- □ *Independency* from vendors of messaging systems
- Coverage of most common facilities that are offered in messaging systems
- It promotes the usage of Java technology

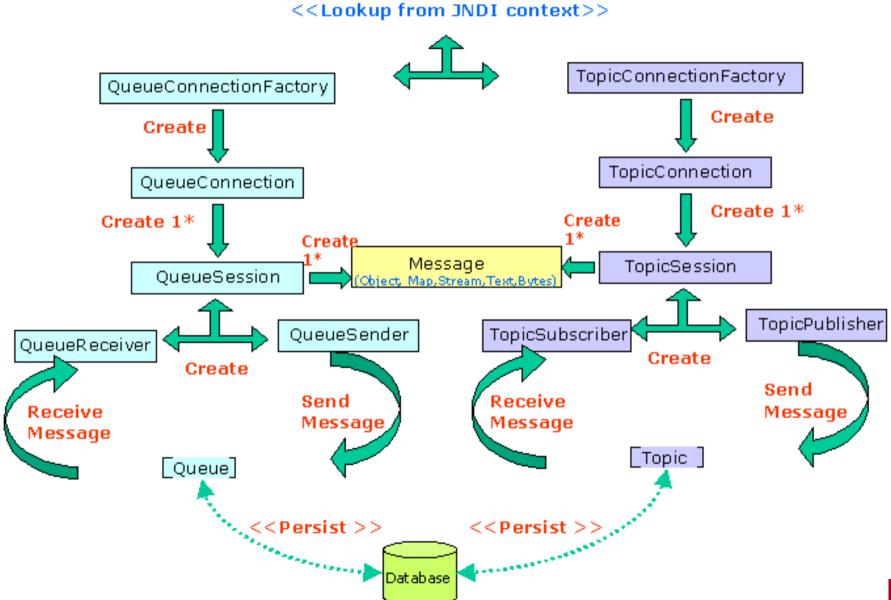
Java™ Application

JMS API



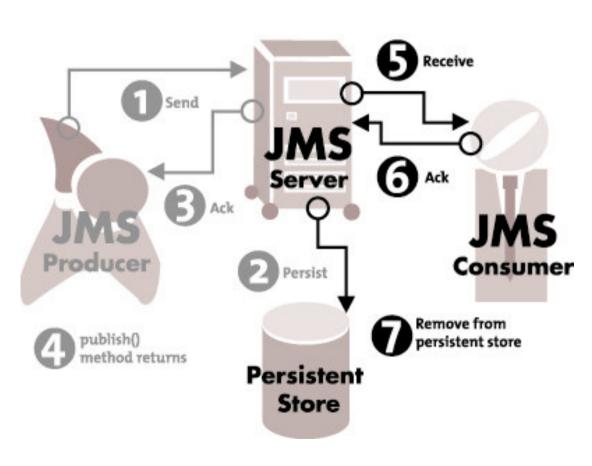


"Graphical Summary" of JMS APIs





Reliability through ACKs: e.g., Automatic ACKs



- Producer-side and consumer-side perspectives
- Differences between persistent and nonpersistent cases
- When is it possible to have duplicated messages?
- When is it possible to have message losses?
- In addition, three differentiated types of ack



Persistency: Two Delivery Modes

□ PERSISTENT

- Default
- It specifies to JMS provider to guarantee that the message is not lost when in transit, e.g., because of a failure at the JMS provider

□ NON_PERSISTENT

- It does NOT request storing messages at the JMS provider side
- Better performance results

<u>SetDeliveryMode()</u> method in the <u>MessageProducer</u> interface

- > producer.setDeliveryMode(DeliveryMode.NON_PERSIS TENT);
- > Extended form: producer.send(message,
 DeliveryMode.NON_PERSISTENT, 3,10000);



Priority and Expiration in Message Delivery

- □ 10 priority levels
 - from 0 (lowest) to 9 (highest)
 - > default = 4

```
Usage of setPriority() method of MessageProducer
interface, e.g., producer.setPriority(7);
  or the extended form producer.send(message,
    DeliveryMode. NON_PERSISTENT, 7, 10000);
```

- Expiration: possibility to configure TTL via setTimeToLive() of the MessageProducer interface
 - > producer.setTimeToLive(60000);
 - > Or extended form, producer.send(message,
 DeliveryMode.NON_PERSISTENT, 3, 60000);



Messaging Examples: CORBA Messaging

CORBA Messaging specification includes:

- □ Asynchronous Messaging Interface (AMI)
 - □ Possibility of both *polling and callback* (callback is passed as CORBA object, therefore even not in the same addressing space of client)
- □ Time Independent Invocation (TII) to specify which
 CORBA objects play the *role of router* for the message
 - Rationale: sender and recipient may be temporarily disconnected
 - □ They compose a *store-and-forward network*

CORBA locator = Interoperable Object Reference (IOR), with different profiles depending on binding protocol

Messages in binary format = Common Data Representation (CDR) Extreme flexibility in the choice of the protocol



CORBA AMI: Callback Mode

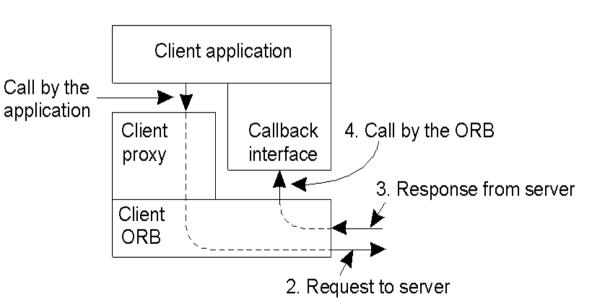
Callback: client provides callback method to be invoked by the support after service completion via a given fire-and-forget (automatically invoked)

In place of: int somma (in int i, in int j, out int somma)
void sendcallback_somma (in int i, in int j)
void callback somma (in int success, in int somma)

Usage of two methods by changing only client implementation and NOT any service part

Client invokes sendcallb...

ORB invokes callback_som..





CORBA AMI: Polling Mode

Asynchronous polling: client decides when and whether to interrogate a method to check completion of remote operation (by collecting results); this method is created by the messaging support

```
In place of: int somma (in int i, in int j, out int somma)
void sendpoll_somma (in int i, in int j)
void pollsomma (out int success, out int somma)
```

Result is collected on request by invoking pollsomma operation that is autom. generated by CORBA support

1. Call by the application

Client proxy | Polling interface | 3. Response from server | Client ORB | 2. Request to server | |



Messaging Examples: Extensible Messaging and Presence Protocol (XMPP)

Essentially designed for instant messaging

RFC 3920 is oriented and similar to the existing implementation of the Jabber protocol; good popularity and widespread utilization thanks to the adoption by *Google, Twitter, Facebook*, ...

It includes *publish/subscribe mechanisms* (see the following slides...) to update presence and state, and for service discovery

Client-server model: client sends an XMPP dataflow to a server, after parameter negotiation

Peer-to-peer model: servers coordinate together for delivery to recipients

Usage of so-called *stanzas*, of three types:

- □ Message stanza one-to-one communication, similar to emails.
- Presence stanza simple pub/sub mechanism, communication is transferred to all subscribers
- □ Info/Query stanza –request-response mechanism



Messaging Examples: Extensible Messaging and Presence Protocol (XMPP)

XMPP messages are streams codified in XML

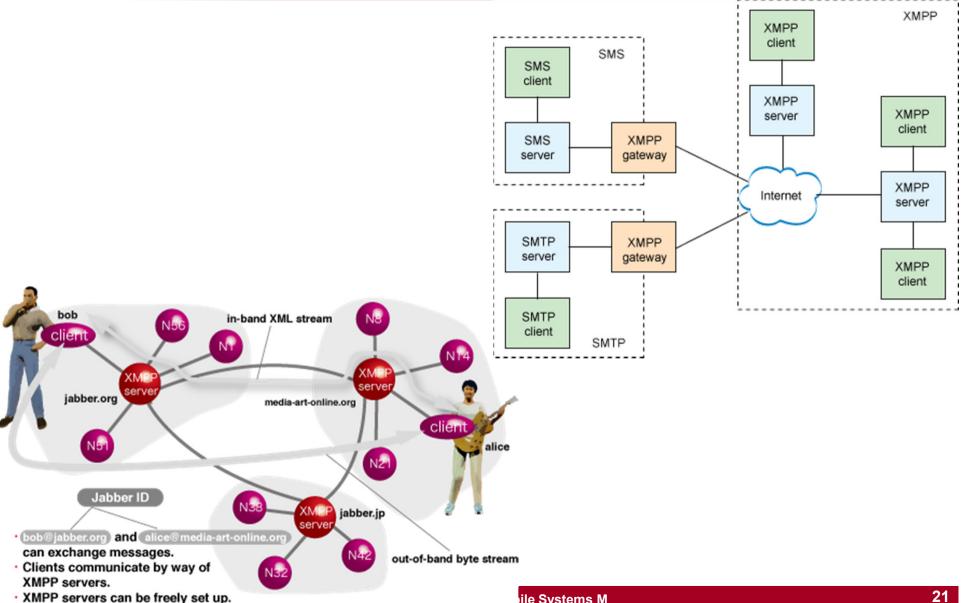
Given the widespread adoption, good candidate to support messaging in mobile systems, EVEN IF:

- Not specifically designed for mobile systems
- Expensive XML processing, expensive connection management in particular in terms of energy
- Expensive re-connections to XMPP server (need to re-establish a new session of interaction per any new transport connection, transmission of XML data that is non-negligible at each session start)

Android implements a specialized and proprietary variant of it, with non-XML-based protocol and NO creation of a new session per any new connection



Messaging Examples: Extensible **Messaging and Presence Protocol (XMPP)**





Messaging Examples: Web Services

SOAP is built on top of interaction model based on message exchange

- Architecture based on senders, receivers, and intermediate nodes
- Locator = HTTP URI
- □ Document-style SOAP: messages as XML-based documents that have to be processed
- □ Possibility of different protocol bindings, but definitely the most used one is HTTP, utilization of POST method (employed more as transport protocol, while ignoring its application semantics)
- In mobile environments, where HTTP is sometimes the only protocol practically usable because of firewalls and NAT, this use/misuse of HTTP could be considered as legitimate and become largely adopted...
- Also specification for binding to email and XMPP



REST is substantially a **solution architectural style**, Resource Oriented Architecture (Roy Fielding, UCI PhD Thesis, 2000)

To promote *client-server* and stateless interaction, oriented to the usage of caching opportunities, also with possibility of code-on-demand to clients

Any resource has a persistent identifier; idea to transfer NOT resources but their representations via HTTP protocol

Constraint: exchange of self-descriptive messages (languages for representation, negotiation of supported modes, ...)



Locator = HTTP URI

Three types of metadata included in HTTP headers:

- Resource metadata about resources, e.g., timestamp about last modification
- Representation metadata about transferred representation, e.g., its media type
- □ **Control metadata** about message, e.g., its length and caching possibility

Notable example: **RESTful Web services**

RESTful Web service as a simple Web service implemented by using HTTP and REST principles, thus resource collection with 3 well-defined aspects:

- □ *URI base for service*, e.g., http://example.com/resources/
- □ *Internet media type* for data used in the service (usually JSON or XML)
- Set of service operations supported via HTTP method invocations (e.g., via POST, GET, PUT or DELETE)



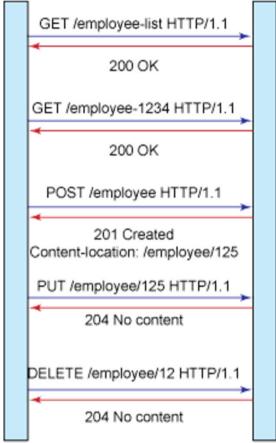
Notable example: **RESTful Web services**

Risorsa	GET	PUT	POST	DELETE
URI for resource collection, e.g., http://example.com/resources/	To list all collection members	To replace the whole collection	To create a new element to be inserted in the collection	To remove the whole collection
URI for single element, e.g., http://example.com/resources/ef7d-xj36p	To obtain the representation of the targeted element, espressed in the appropriate Internet media type	To replace or create an element of the collection	To consider the element as a collection and to create a new element internally to it	To remove an element from the collection

Examples of today's REST usage:

- Majority of Web blogs (download of XML files in RSS/Atom format, which contain links to other resources)
- Simple Storage Service (S3) by Amazon.com
- □ *OpenStreetMap* (REST interface)... and many many others





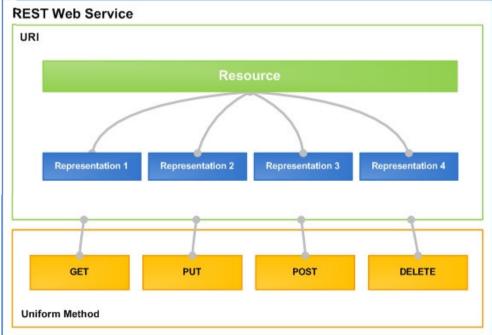
Retrieve a list of all employees

Retrieve a particular employee

Create a new employee and retrieve new ticket URI from HTTP headers

Update an employee

Delete an employee



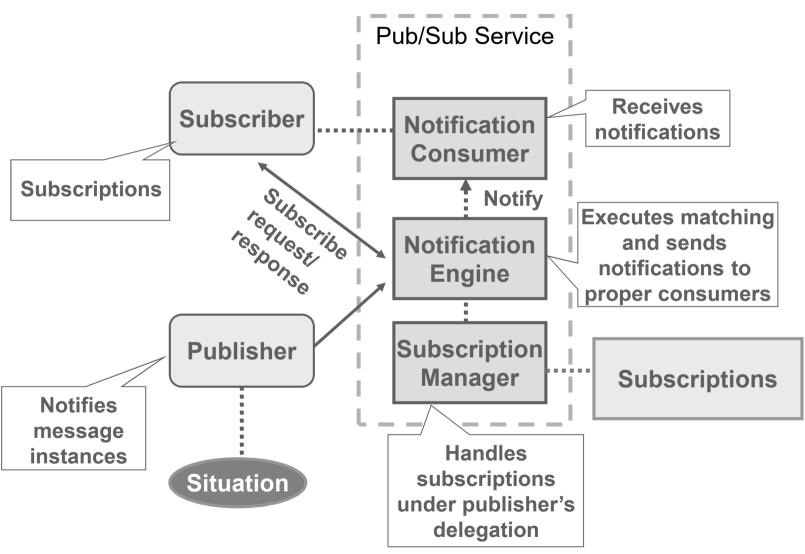


Event Management and Publish/Subscribe Systems

- Event delivery from publishers to subscribers
 - Events as messages with content
 - > **One-to-many, many-to-many** (traditional message systems are queue-based and one-to-one)
 - Often implemented based on messaging systems and on storeand-forward solutions
- □ Comm. paradigm of frequent usage, in particular in mobile systems
 - > Decoupling in space and time
- Event system as logically centralized system
 - Anonymous communication
 - Possibility to use filters (on headers or entire messages)
 - > Basic primitives: subscribe, unsubscribe, publish, also with filters
- Different topologies for routing and different semantics associated to event sending/notification
- Associated operations are typically non-blocking (polling, callback)



General Architecture for Publish/Subscribe Systems





Event Router and Topologies

Event router or broker

- Works as mediator (decoupling) between publishers and subscribers
- Usage of routing table (also with filters) for local event dispatching or to indicate to which «near» router to forward in the case of distributed brokers (to achieve scalability, reliability, and high availability)
- □ Filters may be also based on content => content-based routing
- Other non-functional requirements: notification within time deadlines (bounded delivery time), QoS, fault-tolerance, ordering (causal order, total order)

Possible router topologies:

- Centralized
- Hierarchical (notifications always sent to master, i.e., root of the distribution tree)
- □ **Cyclic, acyclic** (peer-to-peer, cyclic allows redundancy but need of *minimum* spanning tree techniques to prevent from cycles)
- □ Based on *rendez-vous point* (special router that works as rendez-vous, typically for pre-determined types of events)
 - Partially related: have you ever heard of *Distributed Hash Tables (DHTs)*?



Interest Propagation and Subscriptions

One of the primary functions of a router is to *propagate* notifications to near routers that are interested in that event. To this purpose, how to propagate interests and subscriptions?

Properties to be achieved: reduced forwarding overhead, high performance, fast support to variations

- **Simple routing**: any router knows all subscriptions in the global systems (subscription flooding), possibly with optimization of NO forwarding if subscription message has been already circulated
- □ **Covering-based routing**: forwarding of only the more general subscription filters (*which possible issues with unsubscription?*)
- Merging-based routing: it allows to merge different entries in routing table for the sake of table size optimization (usually combined with covering, here also unsubscription issues)

Notifications are usually distributed over *reverse paths (wrt subscription paths)*



Decision about Message Routing

Depending on what is used to take message routing decisions, classification into:

- □ **Channel/topic-based**: depending on the channel (usually named channel) on which the event is published. Pub/sub agreement on the channel name, also possibility of associated multicast address
- □ **Subject-based**: depending on event subject, single field of info
- Header-based: depending on a set of fields. For example, SOAP supports header-based routing for its messages
- □ **Content-based**: possibly depending on the whole message content. Higher expressive power, higher costs

Also *context+content-based routing*, particularly suitable for mobile systems/services with *event filters that are context-dependent*



- Also Java has a built-in *model for event distribution,* based on RMI, e.g., used in Jini/River
- Based on *Remote Event Listener* (consumers are registered to receive given types of events from given objects, **notify()** method)
- □ Remote Event object returned back during notification (data, reference to source object, handback object, unique identifier)
- □ Lease mechanism
- The specification includes possibility to define
 Distributed Event Adaptors that implement filters and QoS policies
 - Idea to exploit handback object, returned by the event source, to transfer state and behavior (e.g., to implement event filters)



```
package net.jini.core.event;

public class RemoteEvent {
    public long getID();
    public long getSequenceNumber();
    public java.rmi.MarshalledObject getRegistrationObject();
}
```

Events generated in local components may transfer even quite complex object state. **NOT distributed events: only info on how state** retrieval is possible at runtime

- Remote event as serializable object that can be transferred between listeners
- □ Idea, "stolen" from Xt Intrinsics and Motif solutions: *to register clients by including handback objects, returned back with any event*

For example, a Jini taxi driver subscribes to taxi bookings while passing through a city area (handback includes location); when it receives an event, it can be informed of old location (at the moment of registration)

Possibility to register other objects for notification delegation: in this case, handback can work as "reminder" with info of subscribers (stock broker model)



Event registration

Jini/River does NOT specify how to register listeners at event sources; only specification to use a class as return value from subscription:

package net.jini.core.event;

Therefore, the developer of event source has to implement:
public EventRegistration
 addRemoteEventListener(RemoteEventListener listener);



Java model for local events work with objects that are all in the same addressing space

Jini as community of distributed objects that cooperate through proxies

For remote events, "inversion of proxy direction"

- For example, Jini client uses its proxy for service access and through it registers itself as listener
 - Need for a proxy method to add event listeners
- Proxy will invoke the "real" method for listener adding over the discovered resource
- □ Invocation of *registration of local event to proxy*; invocation of proxy for remote resource registration

As if real resource obtains a proxy for the client to use in the notification chain



OMG Distributed Data Service (DDS)

- OMG specification (neither based on CORBA nor highly interoperable) for data distribution service designed for real-time systems
- □ Specification defines APIs for so-called publish/subscribe data-centric communication; in other terms, DDS middleware offers abstraction of global data space that is accessible to all interested applications
- Usage of combination of Topic objects and keys to univocally identify instances within a datastream of the same topic
- Support to content filtering and QoS negotiation
- Suitable for distributed propagation of signals, data, and events

CORBA Event Service (NOT data-centric and with NO QoS support);

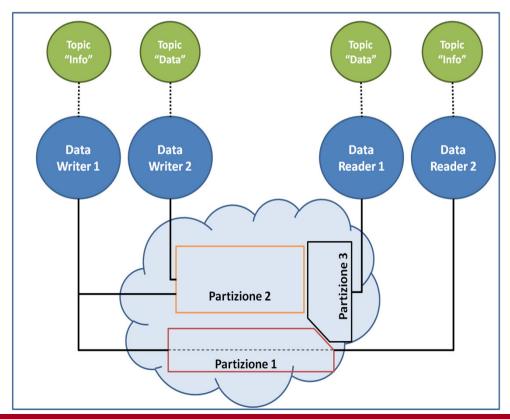
CORBA Notification Service (filters, QoS, but mandatory usage of CDR and IIOP)



Content Subscription: DDS Partitions

Partitions are namespaces to allow the logical splitting of a DDS domain

Publisher/Subscribers can decide *at runtime* (and NOT at instantiation time as for JMS Topics) on which partitions to publish/subscribe data



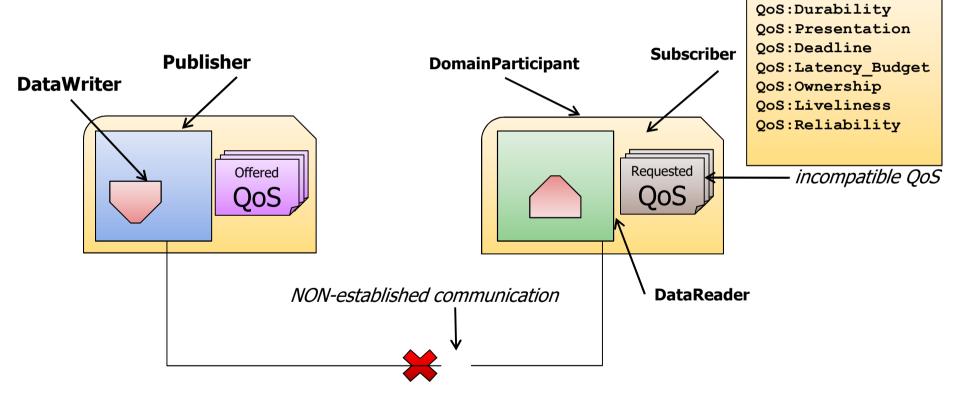
For a DataReader to receive messages from a DataWriter, there is the need to share both the same Topic and the same partition

Partitions are considered to enforce a QoS policy



QoS Negotiation in DDS

- To allow a Subscriber receiving publications from a Publisher,
 QoS properties have to be compatible
- □ Protocol of *Request/Offer negotiation*



DDS supports different modes for message sending (e.g., best-effort, reliable) and personalized management of data persistence



Quality as Reliability

DDS identifies two QoS policies for message reliability:

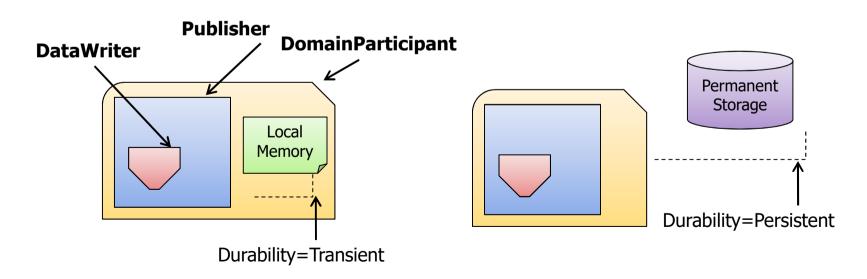
- □ **BEST_EFFORT NOT guaranteed** that all messages are received, NOT guaranteed delivery order
- □ RELIABLE guranteed that all messages are received and delivery order. Via Publishers that re-send data to Subscribers if needed and via Subscribers that send reception feedback (ack)

In reliable case, all sent messages are *kept in a history queue* while waiting for being confirmed (publisher side) and processed by application (subscriber side); queue size can be defined, through *HISTORY policy* It is also possible to define *how many resources* (e.g., memory, max instances) to use to maintain data, through *RESOURCE_LIMITS policy*



Quality as Durability

- Through *Durability policy* it is possible to define whether and how many data to be maintained at publisher side in order to enable their future successive request
- □ DDS supports 3 persistency types:
 - > **VOLATILE** No Instance History Saved
 - TRANSIENT History Saved in Local Memory
 - > **PERSISTENT** History Saved in Permanent storage





DDS Quality: Additional Policies

DDS supports a wide set of other policies to define:

- Ordering of received messages (DESTINATION_ORDER -BY_RECEPTION_TIMESTAMP, BY_SOURCE_TIMESTAMP – eventual consistency, ...)
- Message priority (LATENCY_BUDGET)
- Exclusiveness on some given data types (OWNERSHIP)
- Data authentication and security (USER_DATA)
- □ Time constraints on message sending/delivery rates (TIME_BASED_FILTER)
- □ Fault detection and heartbeat (LIVELINESS)

More detailed technical documents at:

- Getting Started Guide www.rti.com/eval/rtidds44d/RTI_DDS_GettingStarted.pdf
- RTI DDS User's Manual www.dre.vanderbilt.edu/~mxiong/tmp/backup/RTI_DDS_UsersManual.pdf



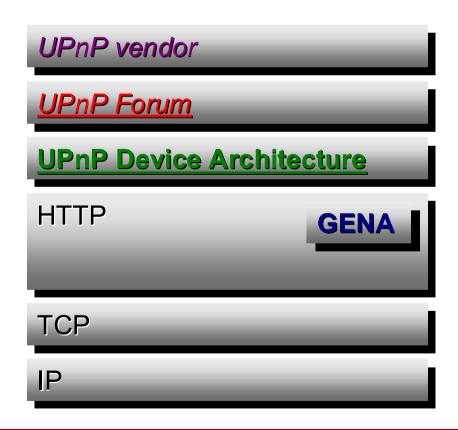
General Event Notification Architecture (GENA)

As already stated, used primarily in UPnP

- Control point is listener of modifications of device state
 - > 0 obtains address
 - > 1 discovers device
 - 2 determines XML descriptor
 - Obtains URL for eventing
 - > 4 registers itself

Extreme simplicity:

Notification sending/reception via HTTP over TCP/IP or multicast UDP





GENA: Subscription

Control point has to register itself before being able to receive any event

SUBSCRIBE publisher path HTTP/1.1

HOST: publisher host:publisher port

CALLBACK: <delivery URL>

NT: upnp:event

TIMEOUT: Second-requested subscription duration

Device accepts subscription: it immediately sends a special event (initial) to control point with the value of all state variables

HTTP/1.1 200 OK

SID: uuid:subscription-UUID

TIMEOUT: Second-actual subscription duration



GENA: Notifications

When a state variable changes value at a device:

```
NOTIFY delivery path HTTP/1.1
HOST: delivery host:delivery port
CONTENT-TYPE: text/xml
NT: upnp:event
NTS: upnp:propchange
SID: uuid: subscription-UUID
SEQ: event key
<e:propertyset xmlns:e="urn:schemas-upnp-org:event-1-0">
 <e:property>
  <variableName>new value/variableName>
 </e:property>
 Other (possible) names of variable and associated values
</<u>e</u>:propertyset>
```



SIP Event Framework

- Asynchronous notifications are essential to implement many SIP services (automated callback services, list of buddies who are currently online, message waiting services, ...)
- To this purpose, *SIP Event Framework* (RFC 3265), based on SUBSCRIBE and NOTIFY methods
- □ SIP events are identified through three elements: Request URI, event type, and message body (optional)
- Notable example: presence service, with so-called presentities and watchers
 - > Presence URI in the format pres:paolo@domain
 - > Scalability issues
 - Need for hierarchical organization into domains and actionf of even aggregation on localities to reduce number of notifications

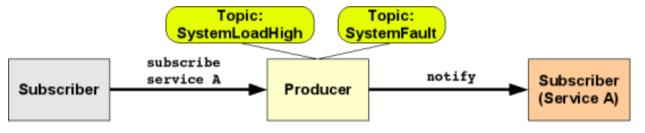


Web Services Event&Notification

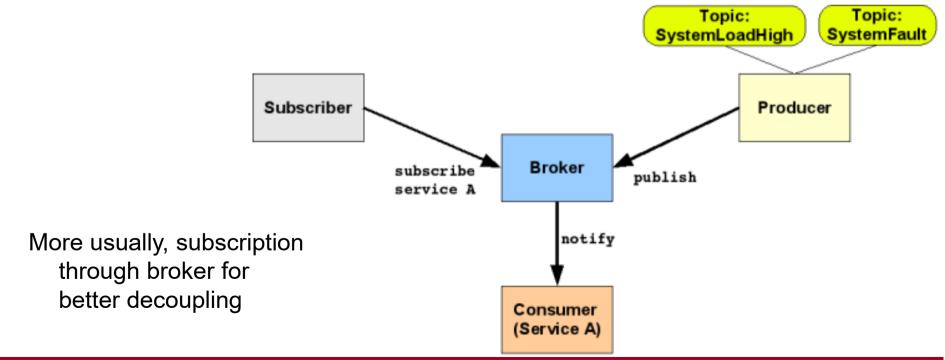
- Two key mechanisms to implement pub/sub for Web services: *WS-Eventing and WS-Notification* (standardization in 2006)
- WS-Eventing is the specification of protocol with which Web services have to make/accept registrations for event notification
 - Mechanisms to create/remove subscriptions
 - Mechanisms to define expiration time and to allow renewal
 - Support to filters (different languages for filter definition may be used)
- WS-Notification is the specification to allow Web services to disseminate data to other Web services
 - Also possibility of organizations oriented to interests (called topics) and interest-based filtering
 - > Distributed topologies for notification brokers



Web Services Event&Notification



Possible subscription from third parties (direct, with NO broker)





Programming Example of WS-Event&Notification

For example, how to implement WS subscriber by using IBM WebSphere:

- □ As usual, need to **obtain WSDL file** for notification broker and subscription manager services (resp. *NotificationBroker.wsdl* and *SubscriptionManager.wsdl*)
- □ If not yet available at client, need to execute wsimport to *generate client stub*
- □ **Look up at notification broker** (need for reference to notification broker service)
- Instantiation of subscription request object and configuration of consumer reference
- Instantiation of subscribe object to include subscription details, like reference to notification consumer

```
import org.oasis_open.docs.wsn.b_2.Subscribe;
import javax.xml.ws.wsaddressing.W3CEndpointReference;
import javax.xml.ws.wsaddressing.W3CEndpointReferenceBuilder;
// Crea oggetto subscription request. DEVE contenere
// ConsumerReference e PUO' includere filtro, InitialTerminationTime
// e SubscriptionPolicy
Subscribe subscribeRequest = new Subscribe();
W3CEndpointReference consumerReference = new
    W3CEndpointReferenceBuilder().address(consumerURI).build();
subscribeRequest.setConsumerReference(consumerReference);
```



Programming Example of WS-Event&Notification

Definition of topic expression as registration filter

It is possible to associate a *Filter object* to registration request to indicate which events are relevant (*filter based on topic, message content, or both*). For example, topic-based filter (with IBM helper classes):

```
import com.ibm.websphere.sib.wsn.jaxb.base.FilterType;
import com.ibm.websphere.sib.wsn.jaxb.base.TopicExpressionType;
// To prepare the topic expression
topicExpression = topicNamespacePrefix + ":" + topicExpression;
TopicExpressionType topicExpressionType = new TopicExpressionType();
topicExpressionType.setExpression(topicExpression);
// To specify mapping from namespace prefix to topic namespace URI
topicExpressionType.addPrefixMapping(topicNamespacePrefix,
  topicNamespace);
// To specify dialect TopicExpression to use
topicExpressionType.setDialect(topicDialect);
// Filter instantiation
FilterType filter = new FilterType();
// To add expression to filter and needed configuration
// subscribe with filter
filter.addTopicExpression(topicExpressionType);
subscribeRequest.setFilter(filter);
```



Programming Example of WS-Event&Notification

Specification of registration duration and request sending

Two modes to specify expiration time for registration:

- 1) namespace URI and Qname objects
- 2) Helper factory = JAXB ObjectFactory

```
import javax.xml.bind.JAXBElement;
import javax.xml.datatype.DatatypeFactory;
import javax.xml.datatype.Duration;
// Option 1: Duration specification (one year from now)
DatatypeFactory factory = DatatypeFactory.newInstance();
Duration duration = factory.newDuration("1Y"; JAXBElement<String>
  initialTerminationTime = new JAXBElement<String>(
     new QName("http://docs.oasis-open.org/wsn/b-2",
     "InitialTerminationTime"), String.class, duration.toString());
// Option 2:
org.oasis open.docs.wsn.b 2.ObjectFactory objectFactory = new org.
  oasis open.docs.wsn.b 2.ObjectFactory();
initialTerminationTime = objectFactory.createSubscribeInitial-
  TerminationTime(duration.toString());
subscribeRequest.setInitialTerminationTime(initialTerminationTime);
org.oasis open.docs.wsn.b 2.SubscribeResponse
  subscribeResponse = port.subscribe(subscribeRequest);
```